Fully Homomorphic Encryption Lab

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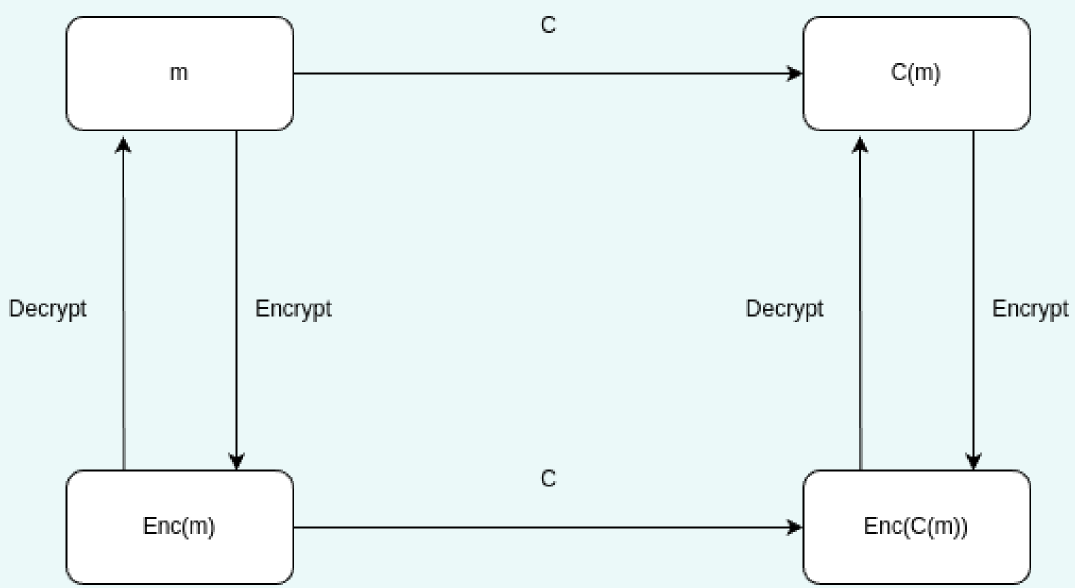

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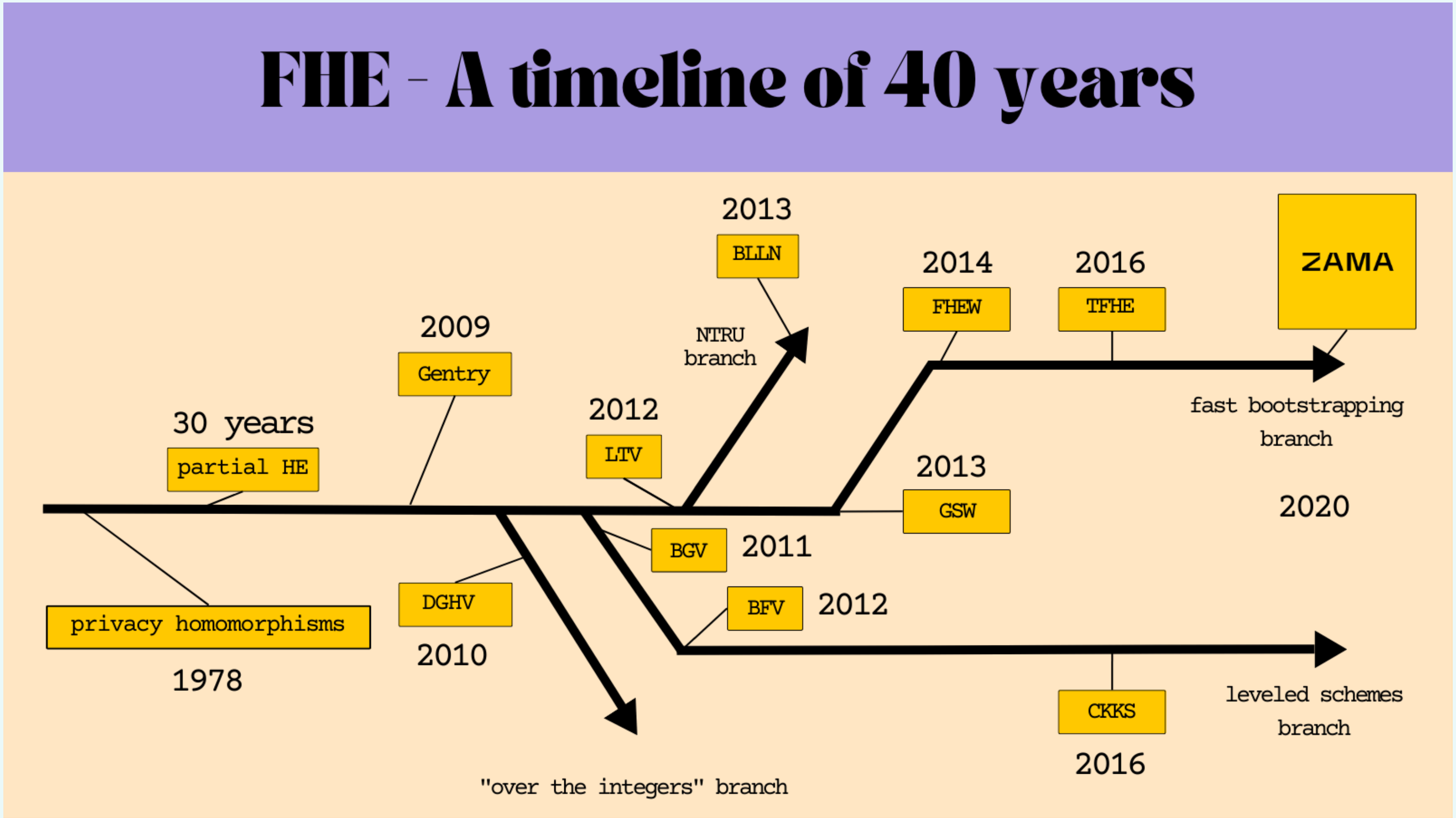
These include:

- Brakerski-Fan/Vercauteren (BFV) [2]
- Brakerski-Gentry-Vaikuntanathan (BGV) [3]
- Cheon-Kim-Kim-Song (CKKS), which supports approximate arithmetic [4]
- Gentry-Sahai-Waters (GSW) [5]
- Homomorphic Encryption over the Torus (TFHE) [6]
- FHEW due to Ducas and Micciancio [7]

A Brief History of ~~Time~~ FHE



Introduction to OpenFHE

OpenFHE is a comprehensive library for employing FHE in code that provides APIs in C/C++, Python and Rust, supporting schemes like BGV, CKKS, and BFV. Here, we demonstrate usage using the Python bindings for OpenFHE's BFV scheme implementation.

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Documentation:

- `https://github.com/openfheorg/openfhe-python`
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Some common methods:

- `MakePackedPlaintext(pt) -> pt1` - Encode the plaintext vector into packed form
- `EvalAdd(ct1, ct2) -> ct3` - Perform homomorphic addition
- `EvalSub(ct1, ct2) -> ct3` - Perform homomorphic subtraction
- `EvalMult(ct1, ct2) -> ct3` - Perform homomorphic multiplication
- `EvalMultKeyGen(sk)` - Generate the relinearization key used for `EvalMult`
- `EvalSum(ct, batchSize) -> ct1` - Evaluate the sum of all components in a vector
- `EvalSquare(ct) -> ct1` - Compute the square of a ciphertext

OpenFHE: Basic usage

This is a sample code template that demonstrates usage of the BFV-RNS (residue number system) scheme for performing FHE (`add` and `mult`) on two plaintext vectors:

```
from openfhe import *

# Set CryptoContext
parameters = CCParamsBFVRNS() # Create instance of the BFV-RNS scheme
parameters.SetPlaintextModulus(65537) # Define plaintext space
parameters.SetMultiplicativeDepth(4) # Max no. of mults w/o bootstrapping

crypto_context = GenCryptoContext(parameters)
crypto_context.Enable(PKESchemeFeature.PKE) # Allow public-key encryption
crypto_context.Enable(PKESchemeFeature.LEVELEDSHE) # Enable leveled FHE w/o
bootstrapping
crypto_context.Enable(PKESchemeFeature.KEYSWITCH) # Enable key switching /
relinearization

# Generate (pk, sk)
key_pair = crypto_context.KeyGen()
```

```python
# Generate the relinearization key
crypto_context.EvalMultKeyGen(key_pair.secretKey)

# Encode first plaintext vector
vec1 = [1, 2, 3, 4, 5, 6, 7, 8, 9, 10]
pt1 = crypto_context.MakePackedPlaintext(vec1)

# Encode second plaintext vector
vec2 = [11, 12, 13, 14, 15, 16, 17, 18, 19, 20]
pt2 = crypto_context.MakePackedPlaintext(vec2)

# Encrypt the two vectors using the same public key
ct1 = crypto_context.Encrypt(key_pair.publicKey, pt1)
ct2 = crypto_context.Encrypt(key_pair.publicKey, pt2)

# Homomorphic addition
ct_add = crypto_context.EvalAdd(ct1, ct2)

# Homomorphic multiplication
ct_mult = crypto_context.EvalMult(ct1, ct2)

# Decrypt the result of the addition
```

```python
pt_add = crypto_context.Decrypt(ct_add, key_pair.secretKey)

# Decrypt the result of the multiplication
pt_mult = crypto_context.Decrypt(ct_mult ,key_pair.secretKey)

print("Plaintext #1: " + str(pt1))
print("Plaintext #2: " + str(pt2))

# Output results
print("#1 + #2 = " + str(pt_add))
print("#1 * #2 = " + str(pt_mult))
```

Output:

```
Plaintext #1: ( 1 2 3 4 5 6 7 8 9 10 ... )
Plaintext #2: ( 11 12 13 14 15 16 17 18 19 20 ... )
#1 + #2 = ( 12 14 16 18 20 22 24 26 28 30 ... )
#1 * #2 = ( 11 24 39 56 75 96 119 144 171 200 ... )
```

Refer to the OpenFHE GitHub repository for more detailed examples that also demonstrate bootstrapping and Threshold-FHE, both of which are beyond the scope of this lab.

Exercises on FHE

You are encouraged to play around with the OpenFHE library! BFV and BGV are two of the simpler FHE schemes with wide-ranging applications. Here are some exercises you can try out:

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- Prove some basic algebraic identities:
  - $(a \pm b)^2 = a^2 \pm 2ab + b^2$
  - $(a^2 - b^2) = (a + b)(a - b)$
  - $(a + b)^3 = a^3 + 3a^2b + 3ab^3 + b^3$
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- Polynomial evaluation:
  - Take a multivariate polynomial, say, $P(x, y) = 2x^2 + 3xy + 4y^2 + 5x + 6y + 7$
  - Evaluate the polynomial homomorphically on, say, $x = 3, y = 4$
  - Decrypt and verify the result $P(3, 4) = 164$

Exercises on FHE (contd.)

- Matrix multiplication:
  - Encrypt two matrices $A = \begin{pmatrix} 1 & 2 \\ 3 & 4 \end{pmatrix}, B = \begin{pmatrix} 5 & 6 \\ 7 & 8 \end{pmatrix}$
  - Perform homomorphic matrix multiplication
  - Decrypt and verify the result matches the product $C = A \cdot B = \begin{pmatrix} 19 & 22 \\ 43 & 50 \end{pmatrix}$

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  - Encrypt a matrix $A = \begin{pmatrix} 1 & 2 \\ 3 & 4 \end{pmatrix}$
  - Compute the determinant $\det(A) = ad - bc$
  - Decrypt and verify the results match $\det(A) = -2$

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- Arithmetic mean:
  - Encrypt a dataset $A = [412, 8423, 66, 891, 277, 84, 5, 9]$
  - Homomorphically compute the arithmetic mean of A
  - Decrypt and verify the result

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